

Authenticating Computation on Groups: New Homomorphic Primitives and Applications

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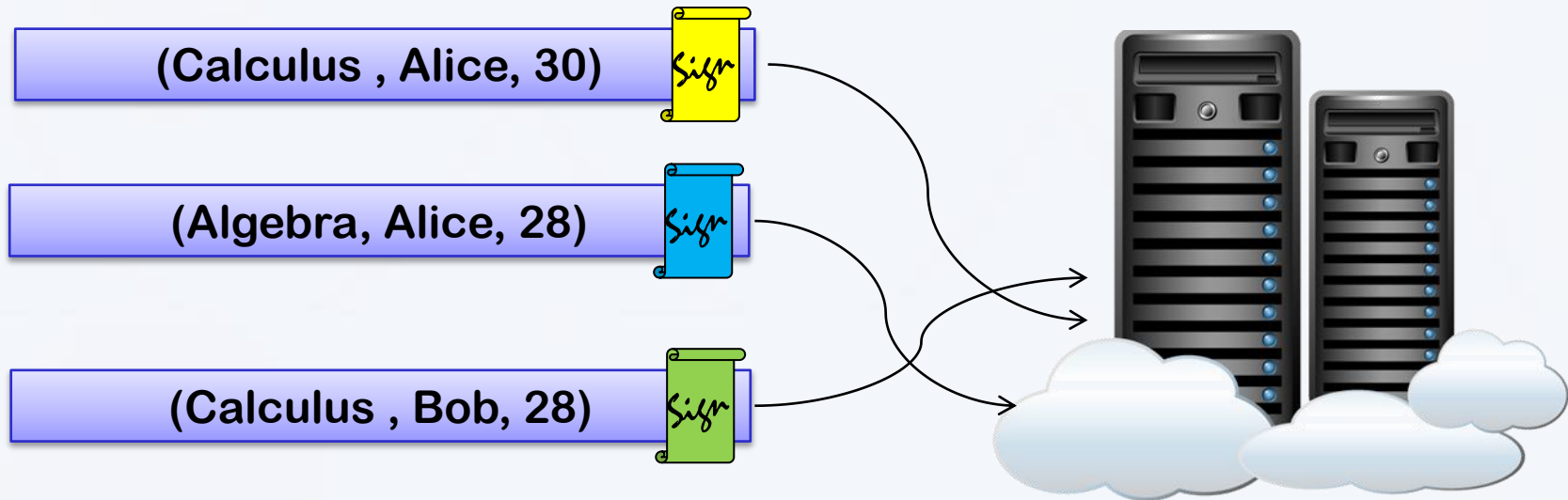
Outline

- ❑ Delegating computation on authenticated data
 - Motivating example
 - Linearly Homomorphic Signature
 - Authenticated Encryption

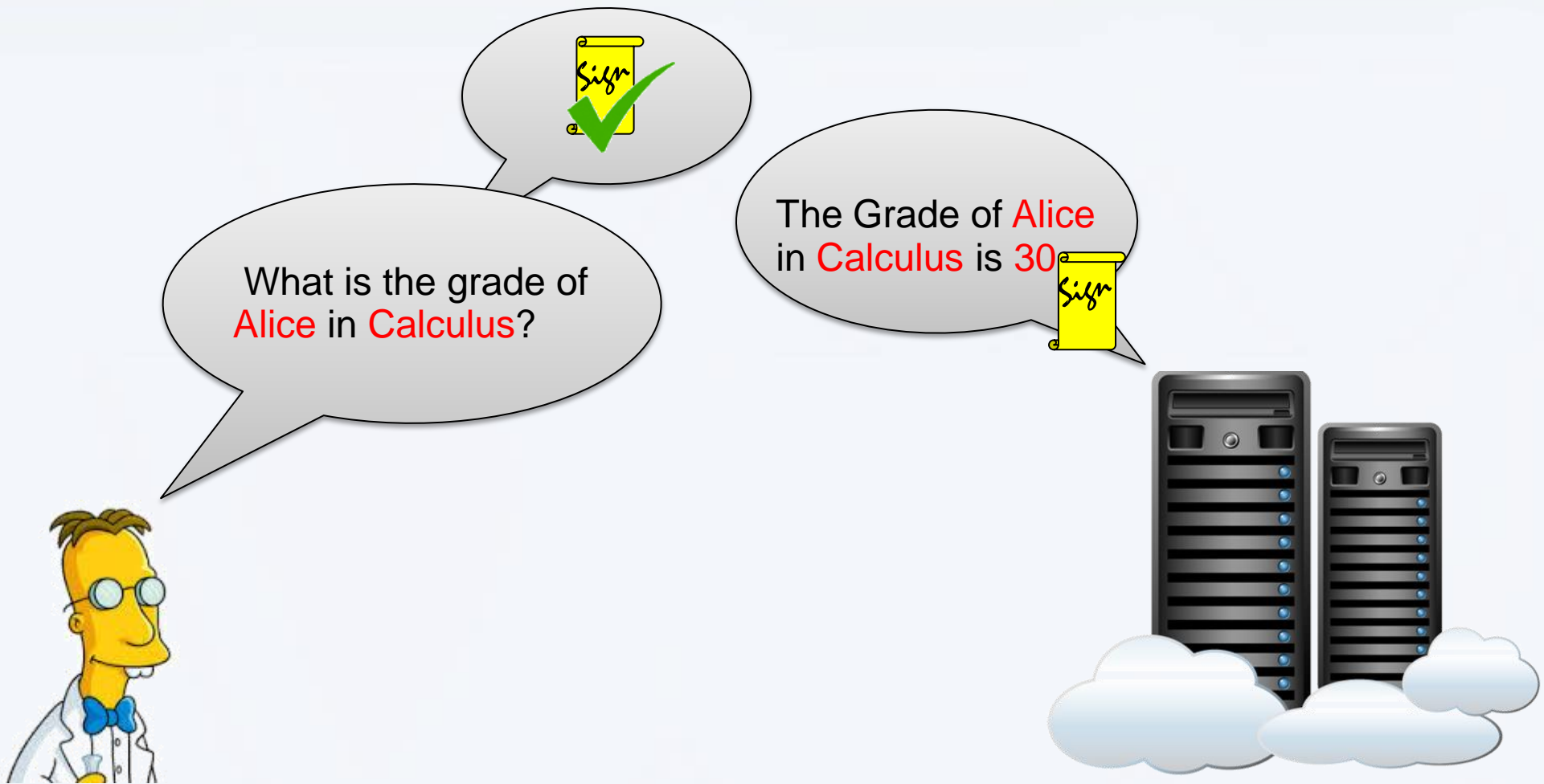
- ❑ Linearly Homomorphic Authenticated Encryption with Public Verifiability (LAEPuV)
 - Definition and security
 - Generic Construction outline and Instantiation

- ❑ Other results

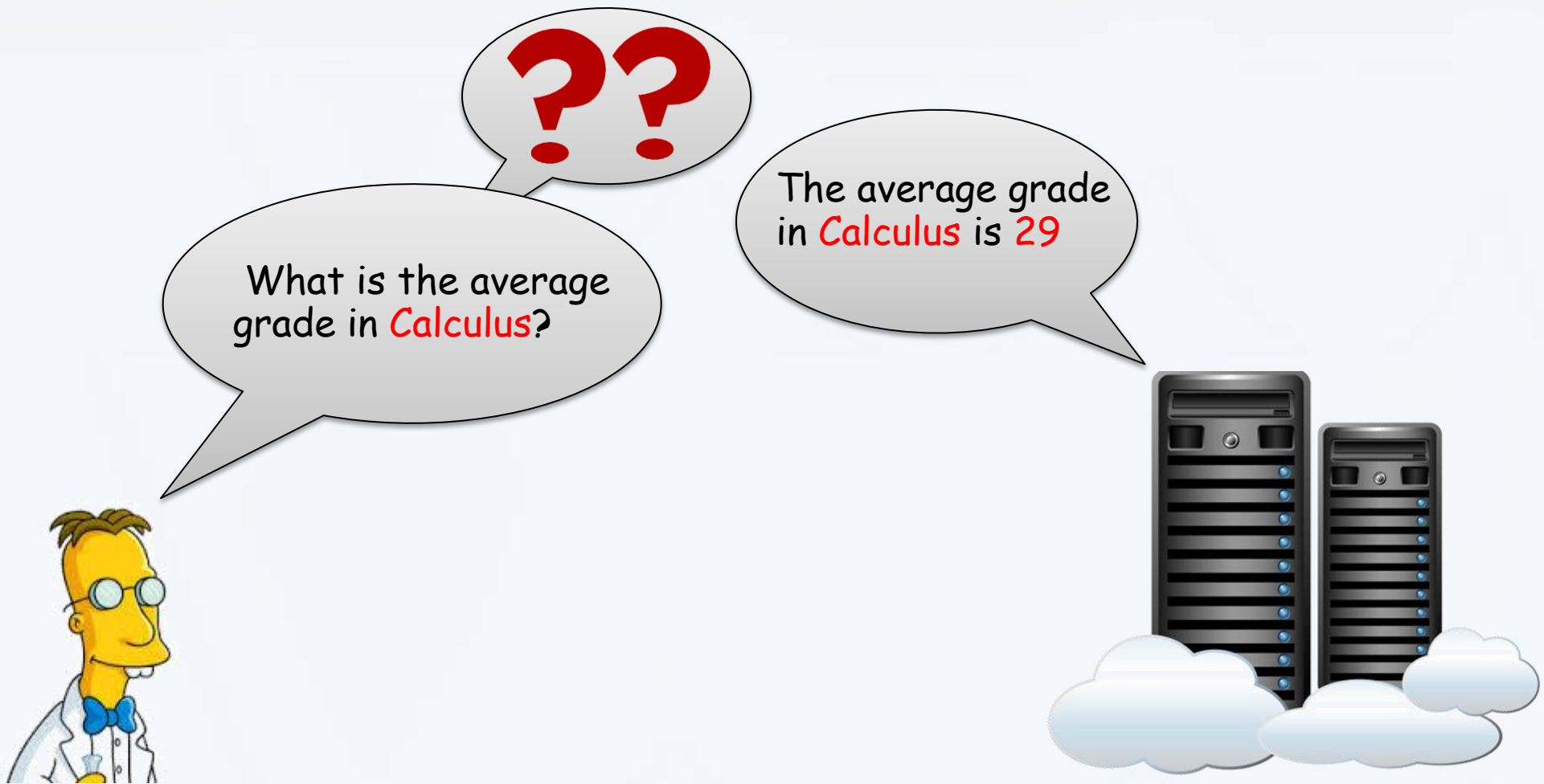
Delegating Computation on Authenticated Data



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Linearly Homomorphic Signatures

[Desmedt93], [JMSW02], [BFKW09]

$(M_1, \text{Sign}(\text{sk}, M_1))$
 $(M_2, \text{Sign}(\text{sk}, M_2))$
.....
 $(M_n, \text{Sign}(\text{sk}, M_n))$



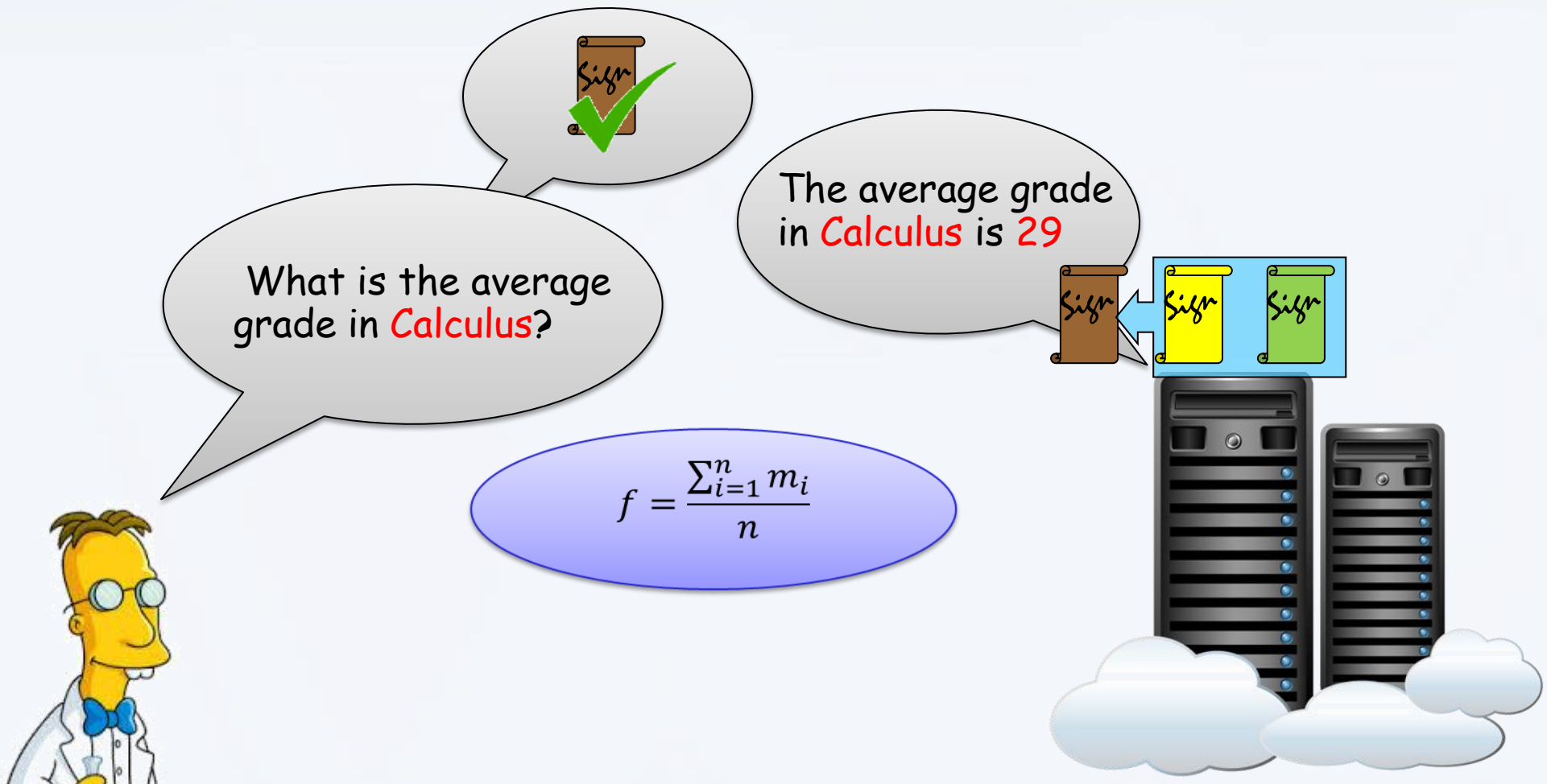
$(f(M_1, \dots, M_n),$
 $\text{Sign}(\text{sk}, f(M_1, \dots, M_n)))$

- Verification is done w.r.t. a function f :

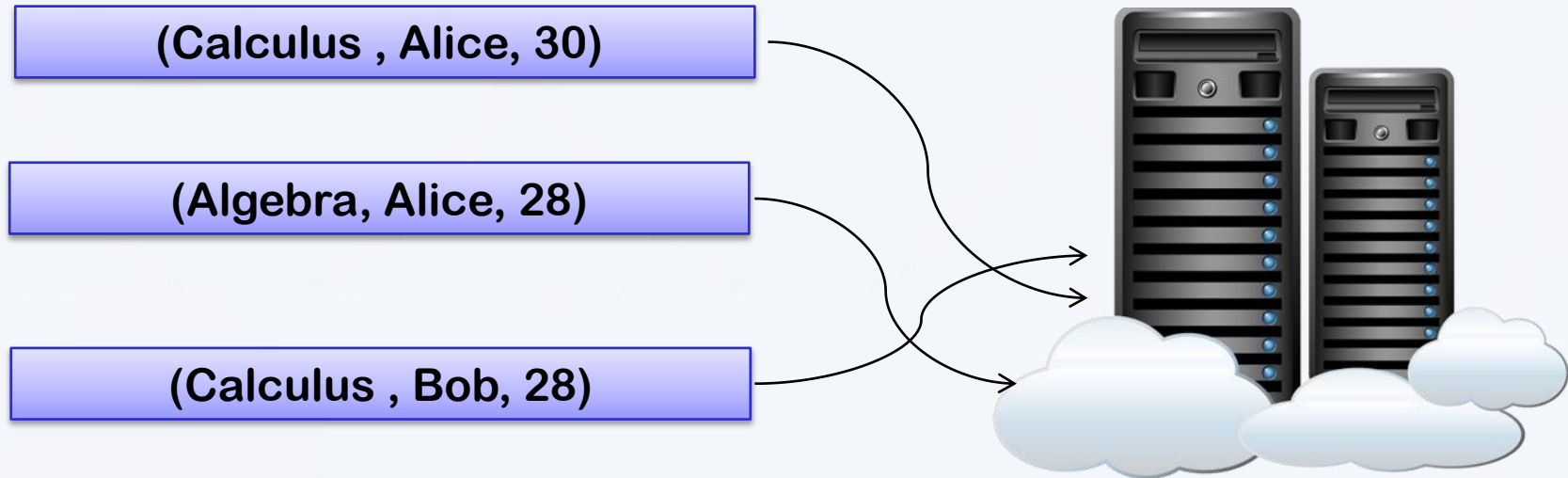
$$\text{Ver}(\text{pk}, f, M, \sigma)$$

- The signatures must be succinct (independent on the number of messages).

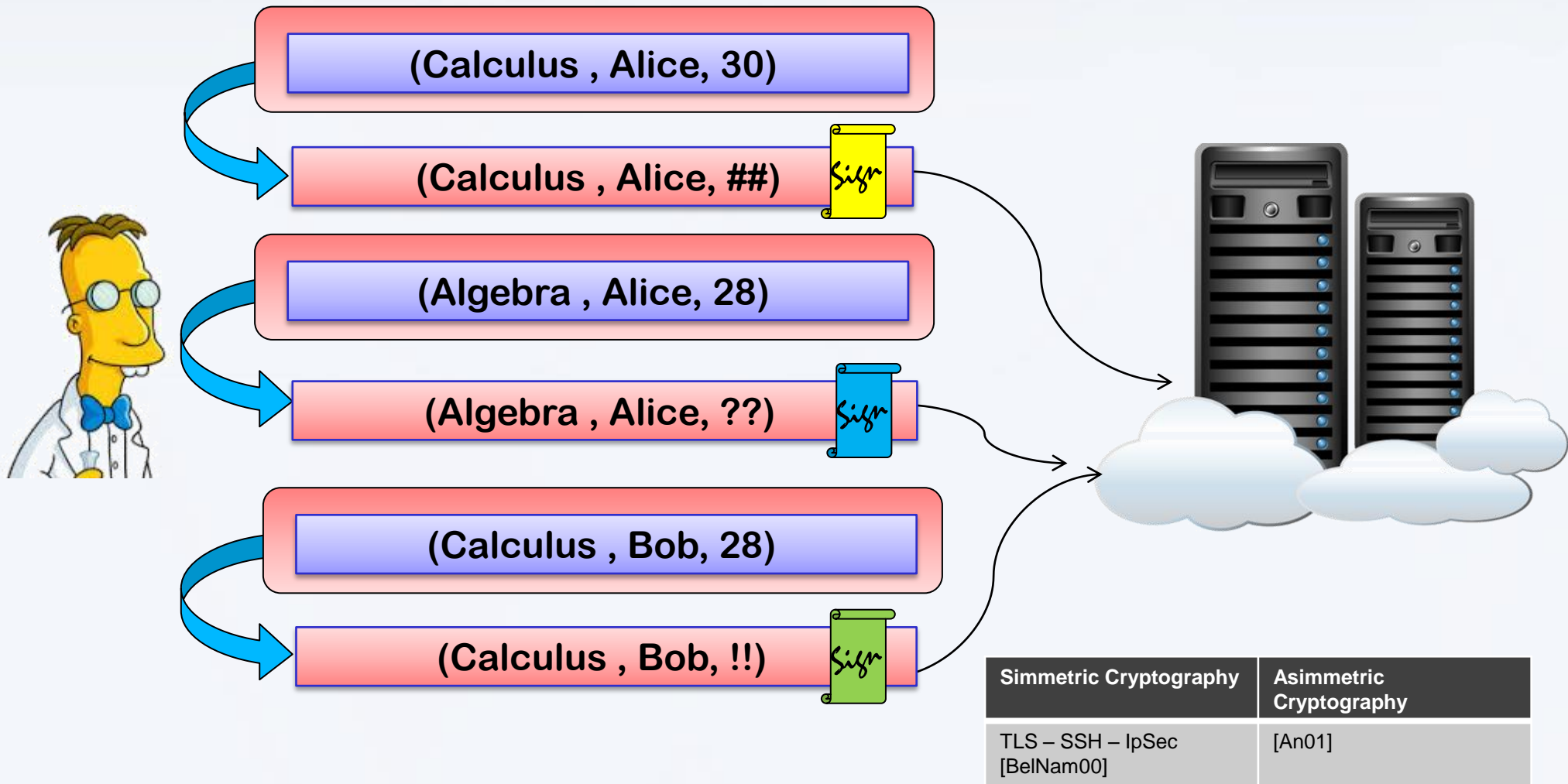
Delegating Computation on Authenticated Data



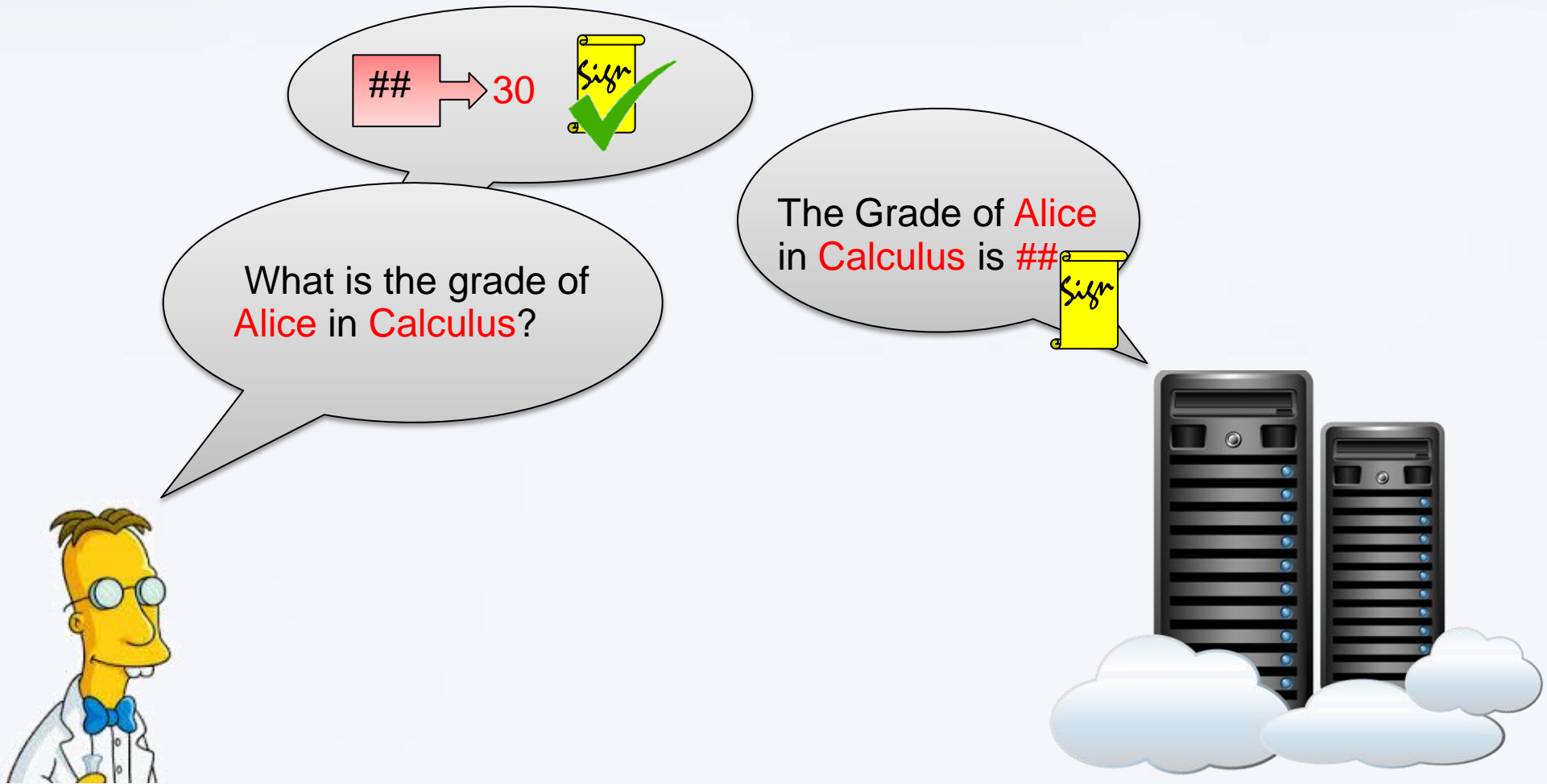
Authenticated Encryption



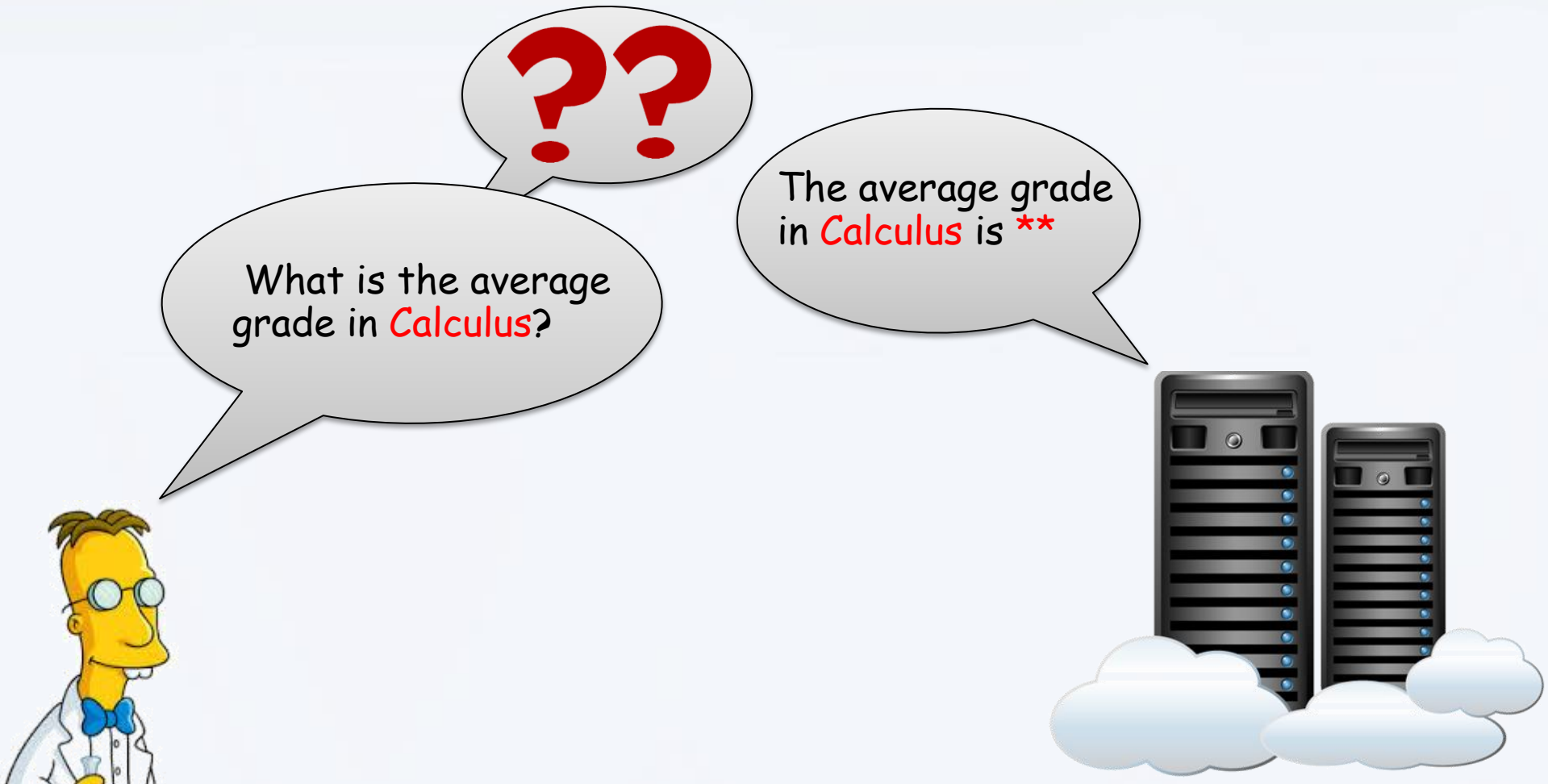
Authenticated Encryption



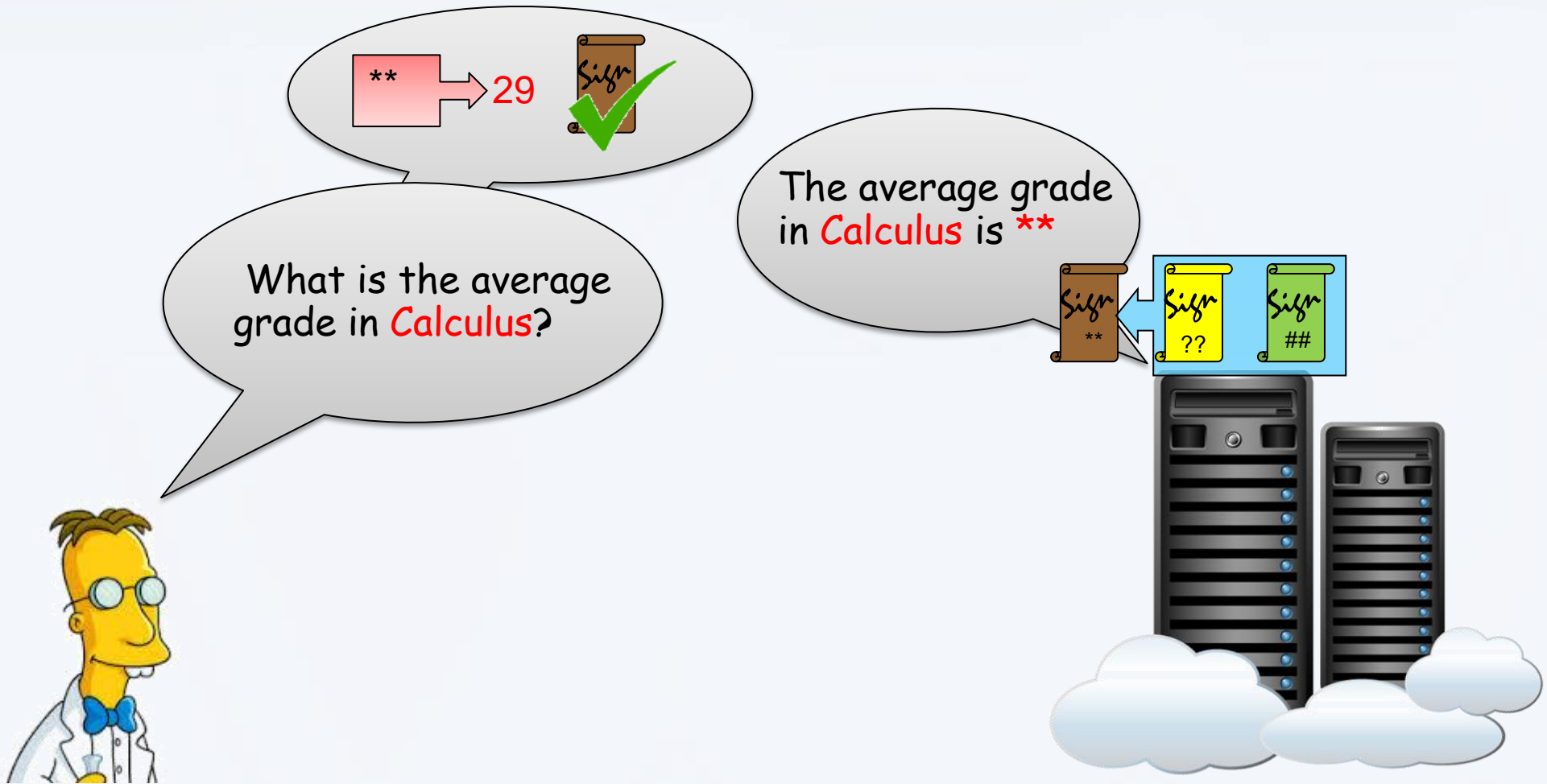
Authenticated Encryption



Authenticated Encryption



Delegating Computation on Authenticated Data with privacy



Linearly Homomorphic Authenticated Encryption with Public Verifiability

Inspired by [JY14]

- $AE\text{-KeyGen}(1^\lambda, k) \rightarrow (sk, vk)$
- $AE\text{-Encrypt}(sk, FID, i, M) \rightarrow C$
- $AE\text{-Verify}(vk, FID, C, f) \rightarrow \{0, 1\}$
- $AE\text{-Decrypt}(sk, FID, C, f) \rightarrow M \text{ or } \perp$
- $AE\text{-Eval}(vk, f, FID, \{C_i\}_{i=1, \dots, k}) \rightarrow C$

Public Verifiability

Security: **LH-IND-CCA** for privacy, **LH-Uf-CMA** for integrity

LAEPuV - General Construction

M message space, **additive** group

R randomness space, **multiplicative** group

C ciphertext space, **multiplicative** group

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T IND-CPA secure

Public Key Encryption Scheme

$$\text{Enc}_{pk}(M_1, R_1) * \text{Enc}_{pk}(M_2, R_2) = \text{Enc}_{pk}(M_1 + M_2, R_1 * R_2)$$

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Σ Linearly Homomorphic signature
scheme for elements in M

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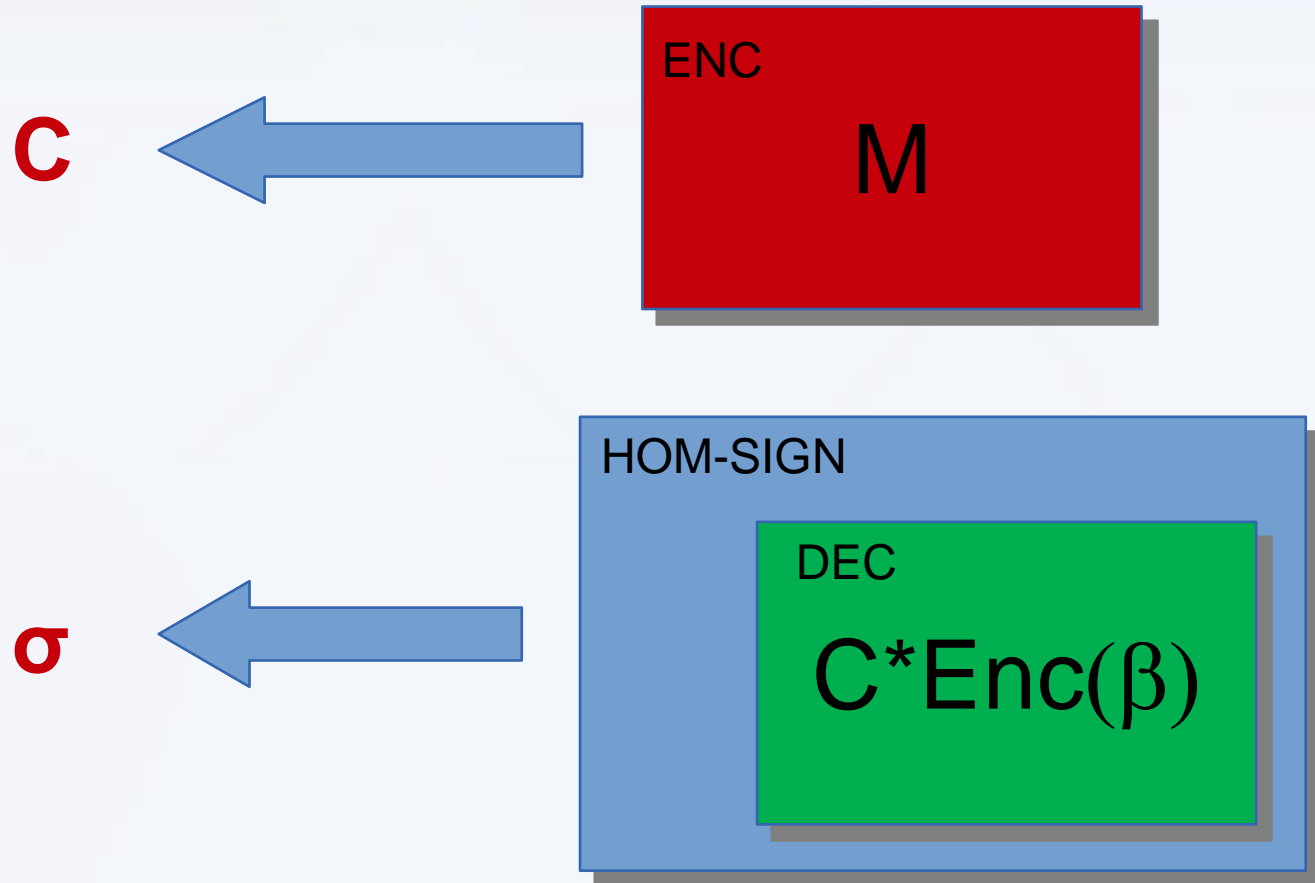
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H Random Oracle

$$H_k: \{0, 1\}^* \rightarrow C$$

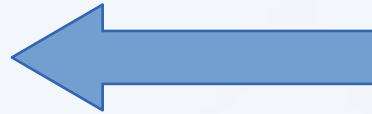
LAEPuV - Encryption



LAEPuV - Eval

$$f = (\alpha_1, \alpha_2, \dots, \alpha_k)$$

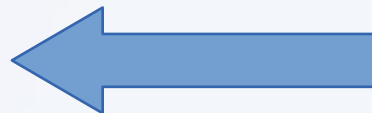
C



$$\prod_{i=1}^k C_i^{\alpha_i}$$

$$f(m_1 + \beta_1, \dots, m_k + \beta_k)$$

σ



HOM-EVAL

f

$$m_1 + \beta_1, \dots, m_k + \beta_k$$

$$\sigma_1, \dots, \sigma_k$$

LAEPuV - Practical Instantiation

- Paillier's encryption guarantees the homomorphic property of \mathbb{T}

$$(g^{m_1} r_1^N)(g^{m_2} r_2^N) = g^{m_1+m_2} (r_1 r_2)^N$$

- As concrete instantiation of the linearly homomorphic signature scheme one can use a simple variant of the (strong) RSA based scheme [CFW12]

Other results (in the paper)

- A Linearly homomorphic signature scheme to sign elements in (bilinear) groups
- This has nice applications in the context of On-line/Off-line signatures

Related works

- **Efficient Delegation of Computation over encrypted data**
 - [JY14], [FGP14]

Conclusion and Open problems

- ✓ Very efficient
- ✓ General construction
- ✓ Public Verifiability


Only linear functions ~~X~~

Needs ROM ~~X~~

Conclusion and Open problems

Interesting Open questions remain :

- How to extend to more general functionalities?
- How to avoid ROM?



Thank You